

# Basis Theorems: Reverse mathematics and Weihrauch reductions

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March 2025

Dagstuhl Seminar 25131

# Basis theorems from reverse mathematics

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Theorem: ( $\text{RCA}_0$ ) The following are equivalent:

- (1)  $\text{I}\Sigma_2^0$ , induction for  $\Sigma_2^0$  formulas.
- (2) Color basis: If  $f : \mathbb{N} \rightarrow k$  then there is a set consisting of those  $j < k$  appearing infinitely often in the range of  $f$ . [1]
- (3) Vector space basis: If  $V$  is a countable vector space of bounded dimension, then  $V$  has a basis. [3]
- (4) e-matroid basis: If  $M$  is an e-matroid of bounded dimension, then  $M$  has a basis. [3]

In their article

*On the Weihrauch degree of the additive Ramsey theorem*

Computability (2024) [4],

Arno Pauli, Cécilia Pradic, and Giovanni Soldà noted that:

many principles equivalent to  $\text{I}\Sigma_2^0$  over  $\text{RCA}_0$  have associated Weihrauch principles related to  $\text{TC}_{\mathbb{N}}^*$ ,  $(\text{LPO}')^*$  or  $\text{TC}_{\mathbb{N}}^* \times (\text{LPO}')^*$ .

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$\text{TC}_{\mathbb{N}}$ : Totalization of closed choice on  $\mathbb{N}$ . Given  $e : \mathbb{N} \rightarrow \mathbb{N}$  choose  $n$  such that either  $n \notin \text{Range}(e)$  or  $\mathbb{N} = \text{Range}(e)$ .

$\text{LPO}'$ : Jump of the limited principle of omniscience. Given infinite binary strings  $\langle p_0, p_1, \dots \rangle$ , apply LPO to  $\lim_i p_i$ .

$\times$  denotes parallel product

$( )^*$  denotes finite parallel products

The previously mentioned basis theorems exhibit this behavior.  
As Weihrauch principles:

- Matroid basis is equivalent to  $\text{TC}_{\mathbb{N}}^* \times (\text{LPO}')^*$  [3]
- Vector space basis is equivalent to  $\text{TC}_{\mathbb{N}}^* \times (\text{LPO}')^*$  [3]
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In general, given

- a bound on the dimension, or
- the exact dimension, or
- a spanning set,

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In general, given

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we would like to find a basis.

How can we reformulate basis problems to modify the Weihrauch strength?



# e-matroid

An e-matroid consists of a countable set  $M$  and an enumeration  $e : \mathbb{N} \rightarrow [M]^{<\mathbb{N}}$  of the finite dependent sets of  $M$ .

Properties:

- $\emptyset$  is independent.
- Finite supersets of dependent sets are dependent.
- If  $X$  and  $Y$  are independent sets and  $|X| < |Y|$ , then for some  $y \in Y$  the set  $X \cup \{y\}$  is independent.

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If  $\text{Range}(e)$  includes all the sets of size  $b$ , then the dimension (rank) of the matroid  $(M, e)$  is bounded by  $b$ .

A basis of a matroid is a maximal independent set.

Let  $\text{EMB}_+$  denote the problem which given input  $(M, e)$  and  $b$ , outputs a basis for  $(M, e)$ .

As noted above  $\text{EMB}_+ \equiv_W \text{TC}_{\mathbb{N}}^* \times (\text{LPO}')^*$  [3]

## Reformulating $\text{EMB}_+$

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Exact dimension reformulation

Let  $\text{EMB}_{\text{rank}}$  denote the problem which given input  $(M, e)$  and the exact rank  $d$  of  $M$ , outputs a basis.

Using a bijection between  $M^d$  and  $\mathbb{N}$  we can find an enumeration of the dependent subsets of  $M$  of size  $d$ . Any set not enumerated is a basis.

Conjecture:  $\text{EMB}_{\text{rank}} \equiv_{\text{W}} \text{C}_{\mathbb{N}}$ .

Conjecture:  $\text{RCA}_0$  proves  $\text{EMB}_{\text{rank}}$ .

Perhaps  $\text{RCA}_0^* \vdash \text{EMB}_{\text{rank}} \leftrightarrow \text{IS}_1^0$ .

## Reformulating $EMB_+$

We should be able to revise  $EMB_+$  to get principles of lower Weihrauch strength.  $EMB_+$  inputs a bound on the dimension.

Spanning set reformulation

Defn: A set  $S$  is a spanning set for a matroid  $M$  if for every  $x \in M - S$  there is an independent set  $I \subset S$  such that  $I \cup \{x\}$  is dependent.

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Let  $\text{EMB}_{\text{span}}$  denote the problem which given input  $(M, e)$  and a finite spanning set  $S$  of  $M$ , outputs a basis.

For each subset of  $S$ , a single use of LPO can determine if the subset is in the range of  $e$ . If it's independent and maximal, then it is a basis.

Conjecture:  $\text{EMB}_{\text{span}} \equiv_W \text{LPO}^*$ .

Conjecture:  $\text{RCA}_0$  proves  $\text{EMB}_{\text{span}}$ .

Perhaps  $\text{RCA}_0^* \vdash \text{EMB}_{\text{span}} \leftrightarrow \text{IS}_1^0$ .

## Color basis recalled

Given  $f : \mathbb{N} \rightarrow k$ , find the color basis, that is, the set of  $j < k$  such that  $j$  appears infinitely often in  $\text{Range}(f)$ .

Theorem:  $(\text{RCA}_0) \mid \Sigma_2^0$  if and only if every  $f : \mathbb{N} \rightarrow k$  has a color basis. [1]

Let  $\text{CB}_+$  denote the problem which given input  $f : \mathbb{N} \rightarrow k$  outputs a color basis for  $f$ .

Theorem:  $\text{CB}_+ \equiv_W (\text{LPO}')^*$  [1]

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Can we reformulate color basis to increase the Weihrauch strength to  $\text{TC}_{\mathbb{N}}^* \times (\text{LPO}')^*$ ?

The color basis is a subset of  $[0, k)$ , so  $k$  is acting like a spanning set.



## Color basis reformulated

Defn: A function  $f : \mathbb{N} \rightarrow \mathbb{N}$  has at most  $b$  color basis elements if, with at most  $b$  exceptions, every element of  $\text{Range}(f)$  appears only finitely many times.

Conjecture:  $(\text{RCA}_0) \mid \Sigma_2^0$  if and only if every  $f : \mathbb{N} \rightarrow \mathbb{N}$  that has at most  $b$  color basis elements has a color basis.

Let  $\text{CB}_{\text{size}}$  denote the problem which given  $b$  and a function  $f : \mathbb{N} \rightarrow \mathbb{N}$  that has at most  $b$  color basis elements, outputs a color basis for  $f$ .

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What about the similar problem with exact size information? Is it related to  $\text{TC}_{\mathbb{N}}^*$ ? Is it equivalent to  $\text{I}\Sigma_2^0$ ?

What about vector space results?

# References

- [1] Caleb Davis, Jeffrey Hirst, Silva Keohulian, Brody Miller, and Jessica Ross, *Reverse mathematics of a pigeonhole basis theorem* (2025). to appear in *Computability*, [Link to a draft](#).
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